

# Reduplication in Tawala isn't Base-Dependent

## Companion Handout

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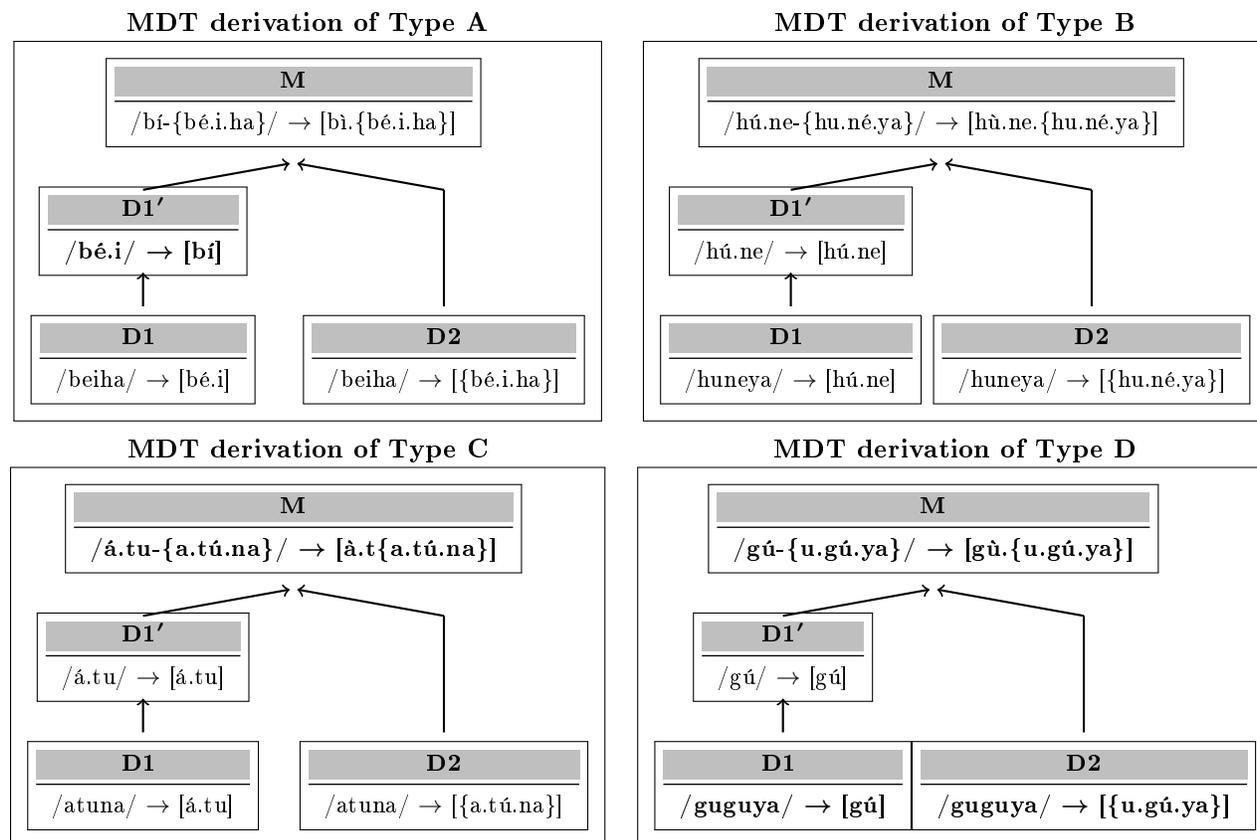
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### 1 Data and Analysis Overview

- Tawala has four distinct, predictable reduplicant shapes:
  - $C_1V_1.V_2$ -initial bases copy  $C_1V_2$ - (**Type A**)
  - CVCV-initial bases copy CVCV- (**Type B**)
  - VC-initial bases copy VC- (**Type C**)
  - Bases that begin in a repeated CV sequence locally double the first root vowel (**Type D**)

	Base	Reduplicated
<b>Type A:</b>	<i>bé.i.ha</i>	→ <i>bì.bé.i.ha</i>
<b>Type B:</b>	<i>hu.né.ya</i>	→ <i>hù.ne.hu.né.ya</i>
<b>Type C:</b>	<i>a.tú.na</i>	→ <i>à.ta.tú.na</i>
<b>Type D:</b>	<i>gu.gú.ya</i>	→ <i>gù.u.gú.ya</i>

- The MDT derivations are given below.
- Main components of analysis:**
  - Truncation in D1
  - Prosodic Root ({...}) parsing in D2
  - Hiatus resolution in M, sometimes blocked by FAITH
  - $V_1$ -deletion under hiatus in D1'



## 2 Truncation in D1

- First step in MDT analysis (IZ:95, HHK): D1 preferentially truncates input down to two syllables (1A,B,C).
  - This is effectuated by the ranking in (2) (constraints defined in (3)), and demonstrated for Type B in (4).
  - It is the combined effect of **STRESSL** and **NONFINALITY** that requires at least two syllables in the truncatum. (This also correctly places stress on the first syllable.)

(1) <b>Output of D1</b> <table style="width: 100%; border-collapse: collapse; border-top: 1px solid black; border-bottom: 1px solid black;"> <thead> <tr> <th style="text-align: left; padding: 2px;">INPUT</th> <th style="text-align: left; padding: 2px;">[D1]</th> </tr> </thead> <tbody> <tr> <td style="padding: 2px;"><b>A</b> <i>beiha</i> → <i>bé.i</i></td> <td></td> </tr> <tr> <td style="padding: 2px;"><b>B</b> <i>huneya</i> → <i>hú.ne</i></td> <td></td> </tr> <tr> <td style="padding: 2px;"><b>C</b> <i>atuna</i> → <i>á.tu</i></td> <td></td> </tr> <tr> <td style="padding: 2px;"><b>D</b> <i>guguya</i> → <i>gú</i></td> <td></td> </tr> </tbody> </table>	INPUT	[D1]	<b>A</b> <i>beiha</i> → <i>bé.i</i>		<b>B</b> <i>huneya</i> → <i>hú.ne</i>		<b>C</b> <i>atuna</i> → <i>á.tu</i>		<b>D</b> <i>guguya</i> → <i>gú</i>		(2) <b>Ranking for 2σ truncation</b> <b>STRESSL, NONFIN</b> ≫ <b>*STRUC</b> ≫ <b>MAX</b>
INPUT	[D1]										
<b>A</b> <i>beiha</i> → <i>bé.i</i>											
<b>B</b> <i>huneya</i> → <i>hú.ne</i>											
<b>C</b> <i>atuna</i> → <i>á.tu</i>											
<b>D</b> <i>guguya</i> → <i>gú</i>											

- |     |    |   |          |
|-----|----|---|----------|
| (3) | a. | <b>STRESSL</b> : Assign one violation if the leftmost syllable is unstressed.               | (*[#σ̥]) |
|     | b. | <b>NONFINALITY</b> : Assign one violation if the rightmost syllable is stressed.            | (*[σ#])  |
|     | c. | <b>*STRUC</b> : Assign one violation for each segment in the output.                        |          |
|     | d. | <b>MAX-IO</b> : Assign one violation for each input segment without an output correspondent |          |

(4) **D1 derivation of Type B (same for Types A & C)**

	/huneya/		STRESSL	NONFIN	*STRUC	MAX-IO
a.	hú.ne.ya	[100]			6!	
b.	hú.ne	[10]			4	2
c.	hú	[1]		*!	2	4
d.	hu	[0]	*!		2	4

- However, Type D, the crucial case in HHK's argument, is truncated to a single syllable.
  - This is effectuated by ranking **\*REPEAT(initial)** (Zukoff 2021; cf. HK) (5) and **STRESSL** above **NONFIN**.
  - This ranking is shown in (6) and demonstrated for Type D in (7).

- (5) **\*REPEAT(initial)**: Assign one violation if the first two syllables are identical. (\*[#σ<sub>α</sub>σ<sub>α</sub>])

- (6) **Ranking for 1σ truncation in Type D**: **STRESSL, \*REPEAT(init)** ≫ **NONFIN**

(7) **D1 derivation of Type D**

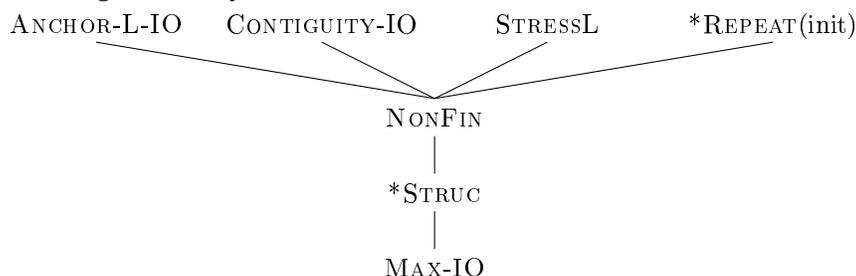
	/gu <sub>1</sub> gu <sub>2</sub> ya/		STRESSL	*REPEAT(init)	NONFIN
a.	gú <sub>1</sub> .gu <sub>2</sub> .ya	[100]		*!	
b.	gú <sub>1</sub> .gu <sub>2</sub>	[10]		*!	
c.	gú <sub>1</sub>	[1]			*
d.	gu <sub>1</sub>	[0]	*!		

- High-ranked **ANCHOR-L** (8a) and **CONTIGUITY** (8b) ensure that other 2σ alternatives are not viable (9).

- |     |    |   |
|-----|----|---|
| (8) | a. | <b>ANCHOR-L-IO</b> : Assign one violation if the leftmost segment in the input does not correspond to the leftmost segment in the output. |
|     | b. | <b>CONTIGUITY-IO</b> : Assign one violation for each pair of adjacent segments in the output which were not adjacent in the input.        |

(9) **D1 derivation of Type D**

/gu <sub>1</sub> gu <sub>2</sub> ya/		ANCHOR-L	CONTIGUITY	NONFIN
a. $\text{gú}_1$	[1]			*
b. $\text{ú}_1.\text{gu}_2$	[10]	*!		
c. $\text{gú}_1.\text{ga}$	[10]		*!	
d. $\text{gú}_1.\text{ya}$	[10]		*!	
e. $\text{gú}_2.\text{ya}$	[10]	*!		

(10) **Ranking summary for D1****3 Type D C<sub>1</sub>-deletion in D2**

- D2 parses all its segment into an output Prosodic Root (cf. IZ:140; Downing 1998a,b).
  - This is effectuated with the constraint MAX-INPUT-PROOT (11), as shown in (12).

- (11) **MAX-IP**: Assign one violation for each input segment which does not have an output correspondent contained within a Prosodic Root.
- (12) **MAX-IP induces PRoot parsing, e.g. for Type A**

/beiha/	MAX-IP
a. bé.i.ha	5!
b. {bé.i}.ha	2!
c. $\{\text{bé.i.ha}\}$	

★ **Note**: it is crucial that D1 not output a PRoot.

- Therefore, D1 needs to have a constraint against PRoot's in the output (e.g. \*PROOT) or against parsing segments into a PRoot (unclear how this would work), which outranks MAX-IP.
- D2 needs to have the reverse ranking: MAX-IP  $\gg$  \*PROOT

- The only constraint which forces violation of MAX-IP in D2 is \*REPEAT(init).
  - If \*REPEAT(init) outranks MAX-IP *and* MAX-IO (13), we generate initial-C deletion for Type D only (14).
  - Since there is not general truncation in D2, we know that \*STRUC ranks below at least one of the MAX constraints.

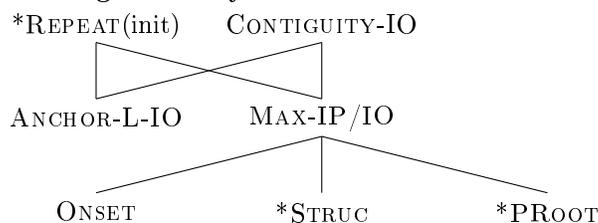
(13) **Ranking for initial-C deletion in Type D: \*REPEAT(init)  $\gg$  MAX-IP, MAX-IO**(14) **\*REPEAT(init)-driven deletion in D2 for Type D**

/gu <sub>1</sub> gu <sub>2</sub> ya/	*REPEAT(init)	MAX-IP	MAX-IO	*STRUC
a. {gu <sub>1</sub> .gú <sub>2</sub> .ya}	*!			6
b. g{u <sub>1</sub> .gú <sub>2</sub> .ya}	*!	*		6
c. $\{\text{u}_1.\text{gú}_2.\text{ya}\}$		*	*	5

- The fact that the \*REPEAT(init) violation is repaired by *initial* deletion and not medial deletion shows that CONTIGUITY-IO  $\gg$  ANCHOR-L-IO (15).
  - At least the higher-ranked MAX constraint must also outrank ONSET, because initial-C deletion (15b) is preferred to initial-CV deletion (15c).
  - The low ranking of ONSET in D2 also explains why hiatus is tolerated stem-internally in D2 (e.g. {bé.i.ha}).

(15) **\*REPEAT(init)-driven deletion in D2 for Type D**

/gu <sub>1</sub> gu <sub>2</sub> ya/	*REPEAT(init)	CONTIG-IO	ANCH-L	MAX-IP/IO	ONSET
a. {gu <sub>1</sub> .gú <sub>2</sub> .ya}	*!				
b. <del>u</del> {u <sub>1</sub> .gú <sub>2</sub> .ya}			*	*	*
c. {gú <sub>2</sub> .ya}			*	**!	
d. {gú <sub>1</sub> .ya}		*!		**	

(16) **Ranking summary for D2**

- D2 also enforces the language's normal stress pattern (see Ezard 1997):
  - Right-to-left trochees (i.e., penult primary stress + alternating stress leftward from the penult).
  - Unless the penult lacks an onset and the antepenult has a higher vowel, in which case stress retracts to the antepenult, as in *bé.i.ha*.
- ★ I will not spell out the stress constraints here, except to note that STRESSL must be low-ranked, unlike in D1.

## 4 Restricted hiatus resolution in M

- The only process in M is deletion to repair hiatus, which occurs at the reduplicant-base juncture for Type C (19).
  - This can be modeled using the ranking ONSET  $\gg$  MAX-IO (18).
  - To prevent deletion of reduplicant-initial vowels in Type C, we could include ANCHOR-L-IO  $\gg$  ONSET in the rankings, or we could use a constraint protecting stressed vowels (17), which will be needed independently below.

(17) **MAX $\acute{V}$ -IO**: Assign one violation for each stressed vowel in the input that lacks a correspondent in the output.

(18) **Ranking for hiatus resolution in M**: MAX $\acute{V}$ -IO  $\gg$  ONSET  $\gg$  MAX-IO<sup>1</sup>

(19) **ONSET-driven reduplicant- $V_2$  deletion in M for Type C**

/á.tu-{a.tú.na}/	MAX $\acute{V}$ -IO	MAX-PO	ONSET	MAX-IO
a. á.tu.{a.tú.na}			**!	
b. á.tu.{tú.na}		*!	*	*
c. <del>a</del> .t{a.tú.na}			*	*
d. t{a.tú.na}	*!			*

<sup>1</sup> If we assume that adjacency relations are established in the input across morpheme boundaries, then CONTIGUITY-IO would have to rank below ONSET as well to allow deletion in Type C.

- Crucially, hiatus is left unresolved in Type D. This can be derived from the combined effect of two special MAX constraints, both of which are sensitive to input properties of segments:
  - A constraint protecting underlyingly stressed vowel: MAX $\acute{V}$ -IO (17).
  - A constraint protecting underlying PRoot segments: MAX-PO (20).

(20) **MAX-PROOT-OUTPUT [MAX-PO]:** Assign one violation for each vowel segment which was part of a PRoot in the input that does not have an output correspondent.

- These constraints, unlike general MAX-IO, rank *above* ONSET (21), and thus can block hiatus resolution.
  - They succeed in doing so across the juncture for Type D (22), because the lefthand vowel is stressed and the righthand vowel belongs to a PRoot.

(21) **Full ranking for hiatus (non-)resolution in M:** MAX $\acute{V}$ -IO, MAX-PO  $\gg$  ONSET  $\gg$  MAX-IO

(22) **Hiatus tolerance across the juncture for Type D in M**

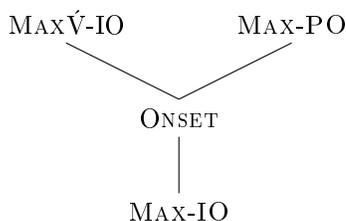
/gú-{u.gú.ya}/	MAX $\acute{V}$ -IO	MAX-PO	ONSET	MAX-IO
a. $\text{gú}\{u.gú.ya\}$			*	
b. $\text{gú}\{gú.ya\}$		*!		*
c. $\text{g}\{u.gú.ya\}$	*!			*

- MAX-PO also explains why there is no hiatus-driven deletion base-internally in, e.g., Type A:

(23) **Hiatus tolerance base-internally for Type A in M**

/bí-{bé.i.ha}/	MAX $\acute{V}$ -IO	MAX-PO	ONSET	MAX-IO
a. $\text{bí}\{bè.i.ha\}$			*	
b. $\text{bí}\{bé.ha\}$		*!		*

(24) **Ranking summary for M**



- The only other piece of phonology in M is demotion of the reduplicant's underlying primary stress to secondary stress.

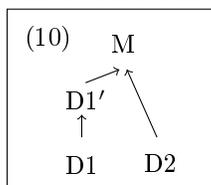
## 5 An extra node for Type A: Hiatus-resolution in D1'

- \* There is one problem remaining with this analysis with respect to Type A (which was suppressed in (23)):
  - Given that the output of D1 for Type A is [bé.i], we predict that the unstressed /i/ should delete under hiatus in M, yielding \*[bè.{bé.i.ha}] (25b).
  - If deletion of the /i/ were blocked by the fact that this would introduce an initial repetition (a \*REPEAT(init) violation), the alternative would be *no deletion* [bè.i.{bé.i.ha}] (25a), not deletion of stressed /é/ (25c,d).
  - And even if something could prefer /é/ deletion, we would still need something else to force the addition of stress to the /i/ in order to overcome (25c).

(25) **Incorrect prediction for Type A in M**

/bé.i-{\b́e.i.ha}/	(*RPT(init))	MAX $\check{V}$ -IO	MAX-PO	ONSET	MAX-IO	DEP[stress]-IO
a. $\text{\textcircled{a}}$ b̀e.i.{b́e.i.ha}				*(!)		
b. $\text{\textcircled{b}}$ b̀e.{b́e.i.ha}	(*!)				*	
c. bi-{\b́e.i.ha}		*!			*	
d. $\text{\textcircled{c}}$ b̀i-{\b́e.i.ha}		*!			*	*

- The simplest solution (following HHK:15–16) is to introduce an additional, semantically vacuous node [D1'] (26) between D1 and M that maps /bé.i/  $\rightarrow$  [b́i] (28).

(26) **MDT derivational structure with D1'**

- The way to do this is as follows:
  - Have ONSET outrank MAX $\check{V}$ -IO and MAX-IO (27a) to make sure stressed vowel deletion will be tolerated as a means of avoiding hiatus.
  - Also have ONSET outrank CONTIGUITY-IO to allow for the discontinuous mapping resulting from V<sub>1</sub>-deletion, and outrank DEP[stress]-IO to allow for restressing the /i/ (27a).
  - Have ANCHOR-R-IO outrank MAX $\check{V}$ -IO and CONTIGUITY-IO (27b) to prefer keeping the rightmost vowel instead of the stressed one, which is also the one that's adjacent to the stem-initial consonant.
  - STRESSL must outrank DEP[stress]-IO (27c) to favor inserting stress onto the newly leftmost /i/, and also outrank NONFIN (27c) (just as in D1) because this means that the final vowel will now be stressed.

(27) **Rankings in D1' to fix Type A**

- ONSET  $\gg$  MAX $\check{V}$ -IO, MAX-IO, CONTIGUITY-IO, DEP[stress]-IO [(28d)  $\succ$  (28a)]
- ANCHOR-R-IO  $\gg$  MAX $\check{V}$ -IO, CONTIGUITY-IO [(28d)  $\succ$  (28b)]
- STRESSL  $\gg$  DEP[stress]-IO, NONFIN [(28d)  $\succ$  (28c)]

(28) **Fixing Type A in D1'**

/bé.i/	ONSET	ANCHOR-R	STRESSL	MAX $\check{V}$	CONTIGUITY	DEP[stress]-IO	NONFIN
a. b̀e.i	*!						
b. b̀e		*!					*
c. b̀i			*!	*	*		
d. $\text{\textcircled{d}}$ b̀i				*	*	*	*

- We'll also need ANCHOR-L-IO to outrank ONSET (29) in order to avoid deleting the stem-initial vowel in Type C (30).

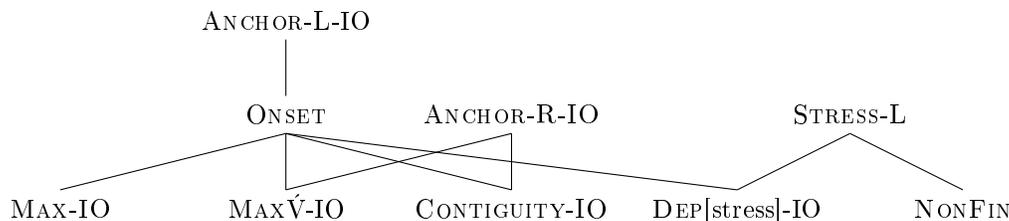
(29) **Ranking for Type C in D1': ANCHOR-L-IO  $\gg$  ONSET**(30) **Not messing up Type C in D1'**

/á.tu/	ANCHOR-L	ONSET	ANCHOR-R	MAX $\check{V}$ -IO	CONTIGUITY
a. $\text{\textcircled{a}}$ á.tu		*			
b. tú	*!			*	

- Once we implement these rankings in D1', we achieve the desired result for Type A in M (31); in fact, there are no obvious competitor candidates.

(31) **Correct result for Type A in M**

/bí- $\{b\acute{e}.i.ha\}$ /	MAX $\acute{V}$ -IO	MAX-PO	ONSET	MAX-IO
a. $\text{b}\acute{i}\text{-}\{b\acute{e}.i.ha\}$				

(32) **Ranking summary for D1'**

## 6 Conclusion

- This analysis, while consistent, requires four substantially different cophologies.
  - The processes that we need to posit in the different nodes bear little resemblance to one another, and sometimes even make contradictory demands.
- MDT provides us with the tools we need to fix all the problems.
  - But the facility with which we can use disparate tools in an ad hoc fashion should give us pause about the restrictiveness of the theory.

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